

2 Number Systems

Student Group

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Table of Contents

2. Number Systems	3
2.1 Types of Number Systems	3
2.1.1 Addition Systems	3
2.1.2 Decimal and Binary	4
2.1.3 Other bases	6
Hexadecimal	6
Octal	8
2.1.4 From decimal to other bases	8
Fast Approach: Try and Subtract	8
Works every time: Repeated Division / Multiplication	8
1. handling the integer decimal	8
2. handling the decimal places	9
2.1.5 Binary Coded Decimals	11
2.1.6 Marking the base of a numeral	11
When to use which base?	11
2.2 basic arithmetic operations in binary and hexadecimal	12
2.2.1 Addition	12
In Binary	12
In Hexadecimal	13
2.2.2 Subtraction	14
In Binary	14
In Hexadecimal	14
Note!	15
2.2.3 Multiplication and Division	15
related Links	16
Exercises	16
Exercise 2.3.1. power tables	16
Exercise 2.3.2. Conversion to Binary	19
Exercise 2.3.3. Conversion to Hexadecimal and Decimal	20

Exercise 2.3.4. Conversion to Hexadecimal and Binary	20
Exercise 2.3.5. Addition and Subtraction	20
Exercise 2.3.6. Multiplication and Division	20
Exercise 2.3.7. Operation in Hex	21
Exercise 2.3.8. Further Questions	21
Exercise 2.3.9. One's Complement	21

2. Number Systems

2.1 Types of Number Systems

In the previous chapter, we had a look at the way a processor (and a computer) can deal with the digital values '\$0\$' and '\$1\$'. However, we haven't seen how the processor can handle larger numbers. To approach this, first a short historical outline is shown.

2.1.1 Addition Systems

The first used number system was the **addition system**. These are also still in use: when enjoying a German beer in the beer garden the waiter is counting the 'progress' by dropping dashes onto the coaster (see [figure 1](#))

Fig. 1: number of drinks on a German Bierdeckel (coaster)



In ancient Rome, these systems were deeper elaborated. Different symbols represent numbers of various sizes:

- $\text{I} = 1$
- $\text{V} = 5$
- $\text{X} = 10$
- $\text{L} = 50$
- $\text{C} = 100$
- $\text{D} = 500$
- $\text{M} = 1000$

Besides this representation of quantities also the position of the symbol in the **numeral** was important:

- In general: the letters have to be arranged decreasing from left to right. For example $\text{MDCI} = 1601$
- There are deviations of this rule: When up to three of the lower symbols are written to the left, these have to be subtracted. Sounds complicated?
A kind of.... for example $\text{MCCD} = 1600$
 $\text{MCCDL} = 1601$
Luckily, we do not have to learn this, but by this trick, the length of the numeral could be shortened,

It becomes even more complicated when trying to calculate with the numbers: what is the result of the multiplication $\text{CCMXXXVII} \cdot \text{DDIIX}$?

2.1.2 Decimal and Binary

Luckily, we have nowadays a better system for writing numbers: the [positional system](#).

We 'just know' what a number like 23 means. However, for understanding how the computer works we have to investigate this gut feeling and put some technical terms onto it.

1. We are accustomed to counting with our fingers from 1 to 10 . For this, we have 10 symbols to count: $0, 1, 2, 3, 4, 5, 6, 7, 8, 9$. This group of distinguishable symbols is called **digits**.
2. The amount of the digits is called **base** B . We are used to the decimal base $B=10$, in logic we used binary (also called dual) $B=2$.
3. When we count beyond the maximum number we are used to 'enlarge the number to the left': after the 9 , we count 10 . But the digit 1 in 10 is more worth than the 1 in 31 (of course). It is on a different **position**, on the position of the tens.
4. Each position gets numbered: the ones count 0, the tens 1, the hundreds 2, the thousands 3, and so on. This 'position number' is called **index** i .
5. Knowing the index, also the 'worth of the position' can be derived: the **place factor** p (like one, ten, hundred) can be calculated with the base and the index: $p=B^i$.
6. A **numeral** as a group of digits represents what is commonly known as a number.
7. A **code** or **encoding** means a way to translate one way to display information into another. E.g. A decimal numeral into a Binary, or an idea of an algorithm into a computer language.

To recapitulate this for $B=10$ we will calculate the amount of a decimal numeral here once in detail (click on the arrow to the right ">" to see the next step, alternatively see [here](#)):

Ok, that was simple. But what about a binary number? Let's calculate the amount of a binary numeral:

So, what did we find out?

- The shown process is a relatively simple way to convert binary numerals to decimals.
- A 8-digit binary numeral is equal to 3 digit decimal numeral. $\rightarrow$ Numerals in binary become lengthy

Therefore, it would be better to have a more structured way of presenting the numerals which are used in the processor. Internally, the processor just knows 0's and 1's. But investigating a huge bunch of these (e.g. when analyzing the internal memory or a file) is not catchy in order to understand anything.

The first step is to group the bits:

- 4 bits are called a **nibble** (the name derives from 'to bit' and 'to nibble')
- 8 bits are called a **byte**
- 16 bits are (usually) called a **word**. In detail, this depends on the processor.
- 32 bits are (usually) called a **double word**. Like the word this also depends on the processor.

By this, one can separate parts of information (e.g. in a file) better. It is also important to mark the order of the bits. For decimal numerals like \$42\$ the rightmost digit has always the lowest value. The technical term for the 'lowest value' in a binary numeral is called **lowest significant bit** or **LSB**. When the LSB is on position 0 (= rightmost) this order is called **LSB 0**. This was used in the calculation above and is commonly used. In some cases (some memory setup and communication protocols) the order is just the other way around. In this case, the **most significant bit** is on position 0. This order is therefore called **MSB 0**. Example: $\$3_{10} = 0000;0011_{2} \text{ (LSB 0)} \$$, $\$3_{10} = 1100;0000_{2} \text{ (MSB 0)} \$$

What is still missing are expressions for large amounts of data. We can describe these using prefixes and the powers of two. You may already know the prefixes such as kilo and mega, but it is worth brushing up on the powers of two. The easiest way to illustrate this is with a chessboard. Maybe you know the legend of the inventor of the chessboard, who was granted a wish by the king. His wish was that the king would give him one grain of rice on the first chessboard square and every time twice as much on each subsequent square. We'll play through this briefly, here to write down the powers of two:

- in the first square we enter two to the power of 0, which results in 1.
- In the second square, two to the power of 1, resulting in two.
- Then 4, 8, 16, 32, 64, 128, 256, 512, and then in the next line two to the power of ten, resulting in 1024.

You should definitely remember this sequence of the values for the power of 2! They are not only important for the exam, but also for the following semesters and computer science.

Fig. 2: chessboard with doubling numerals



1024 bit is also called **kbit** or **kilobit** in the semiconductor industry. You will notice here that the kilo is slightly more than 1000. To make it easier to distinguish, according to the ISO or ECE standard, you should say **kibi** instead of "kilo" and write kibibit. The same applies to 2^{20} , i.e. 1024 to the power of two: **megabit** has become common there. However, **mebibit**, should be used to differentiate. The nomenclature continues in the same way for 2^{30} and 2^{40} : gibibit and tebibit.

2.1.3 Other bases

When looking at multiple bases, it is important to clearly **mark the base** of the numerals. Already in the chapter before a numeral like \$110\$ could either be \$110_{10}\$ in decimal or \$110_2\$ in binary, which is \$6_{10}\$ in decimal.

In the following the base will be written as a subscript: \$110_2 = 6_{10}\$. At the end of this subchapter, we will also see other ways to mark the base.

Hexadecimal

With the ideas of the previous subchapter, we already can structure the bits. In order not to use the lengthy binary presentation it is common to use other bases. One important base is \$16_{10}\$ for hexadecimal numerals. There, we need 16 distinguishable symbols: \$0, 1, 2, 3, 4, 5, 6, 7, 8, 9, \text{A}, \text{B}, \text{C}, \text{D}, \text{E}, \text{F}\$. With these digits, it is possible to encode (=rewrite by other means) exactly 4 bits or one nibble. The encoding will be as follows:

Dual	Decimal	Hexadecimal
\$0000_2\$	\$0_{10}\$	\$0_{16}\$
\$0001_2\$	\$1_{10}\$	\$1_{16}\$
\$0010_2\$	\$2_{10}\$	\$2_{16}\$
\$0011_2\$	\$3_{10}\$	\$3_{16}\$
\$0100_2\$	\$4_{10}\$	\$4_{16}\$
\$0101_2\$	\$5_{10}\$	\$5_{16}\$
\$0110_2\$	\$6_{10}\$	\$6_{16}\$
\$0111_2\$	\$7_{10}\$	\$7_{16}\$
\$1000_2\$	\$8_{10}\$	\$8_{16}\$
\$1001_2\$	\$9_{10}\$	\$9_{16}\$
\$1010_2\$	\$10_{10}\$	\$\text{A}_{16}\$
\$1011_2\$	\$11_{10}\$	\$\text{B}_{16}\$
\$1100_2\$	\$12_{10}\$	\$\text{C}_{16}\$
\$1101_2\$	\$13_{10}\$	\$\text{D}_{16}\$
\$1110_2\$	\$14_{10}\$	\$\text{E}_{16}\$
\$1111_2\$	\$15_{10}\$	\$\text{F}_{16}\$

Tab. 1: hex encoding

This directly reduces the necessary amount of digits to show data. The hexadecimal representation is for example used in the file type *.hex. This is the output file of an embedded c compiler and contains code in a machine-readable representation. An example of a c code and its hex file is shown in [figure 3](#). In the hex file the bytes (= 2 nibbles = 2 digits) are visibly grouped.

```

41 // Deklarationen -----
42 // Festlegung der Quartzfrequenz
43
44 #ifndef F_CPU
45 #define F_CPU 1843200UL // optional definieren
46 #define F_CPU 1843200UL // MinIMEXXL mit 18,432 MHz Quartz
47 #endif
48
49 // Include von Header-Dateien
50
51 #include <avr/io.h> // I/O Konfiguration (intern weitere Dateien)
52 #include <util/delay.h> // Definition von Delays (Wartezeiten)
53 #include "lcd_lib_de.h" // Funktionsbibliothek zum LCD-Display
54
55 // Konstanten
56
57 #define MIN_PER 100 // minimale Periodendauer in "Timerticks"
58 #define MAX_PER 255 // maximale Periodendauer in "Timerticks"
59 #define WAIT_TIME 2000 // Wartezeit zwischen Taktungen in ms
60
61 // Funktionsprototypen
62
63 void initdisplay(void); // Initialisierung Display und Startanzeige
64 void initPorts(void); // Initialisierung der I/O-Ports
65 void initTimer(void); // Timer 0 initialisieren (Soundgenerierung)
66 void init(void); // generelle Initialisierungsfunktion
67
68 // Hauptprogramm -----
69
70 int main()
71 {
72     init(); // Ports und Timer 0 initialisieren
73     initdisplay(); // Display aktivieren
74
75     while(1) // Start der unendlichen Schleife
76     {
77         for (OCR0A=MAX_PER; OCR0A>MIN_PER; OCR0A--) // Frequenz erhöhen
78         {
79             _delay_ms(10); // in Schritten von 10ms
80         }
81         _delay_ms(WAIT_TIME); // Wartezeit hohe Frequenz
82
83         for (OCR0A=MIN_PER; OCR0A<MAX_PER; OCR0A++) // Frequenz absenken
84         {
85             _delay_ms(10); // in Schritten von 10 ms
86         }
87         _delay_ms(WAIT_TIME); // Wartezeit niedrige Frequenz
88     }
89 }
90 // Ende der unendlichen Schleife

```

Fig. 3: example of c code and hex-file

```

1 :1000000019C02BC02AC029C028C027C026C025C0BF
2 :1000100024C023C022C021C020C01FC01EC01DC0DC
3 :100020001CC01BC01AC019C018C017C016C015C00C
4 :1000300014C013C011241FBECFEFD4E0DEBFCDBF6C
5 :1000400011E0A0E0B1E0EAE1F2E002C005900D921B
6 :10005000A433B107D9F7B6D0DEC0D2CF5F9A5F988C
7 :100060000895CF93C82F88B1807F88B988B19C2F1D
8 :1000700092959F70892B88B9F1DF88B1807F88B90C
9 :1000800088B1CF70C82BC8B9E9DFCF910895F89433
10 :100090009FE097B98AB180698AB998B98BB18F7696
11 :1000A0008BB92FEF8FE191E0215080409040E1F734
12 :1000B0000C0000088B1807F88B988B1836088B9AA
13 :1000C000CDDF8FEF93EB0197F1F700C0000C6DFA3
14 :1000D0008FEF93EB0197F1F700C00000BFDF88B10D
15 :1000E000817088B98FEF93E20197F1F700C00000AB
16 :1000F000B5DF86EE90E00197F1F7000088E2B1DF0E
17 :1001000096EF9A95F1F78CE0ACDF86EE90E00197E0
18 :10011000F1F7000082E0A5DF86EE90E00197F1F7AD
19 :10012000000081E09EDF8FEF95E30197F1F700C0BB
20 :10013000000081E096DF8FEF95E30197F1F700C0B3
21 :100140000000789408955C9A8CDF8DEF90E0019721
22 :10015000F1F700000895CF93DF93EC0188818823A5
23 :1001600029F02196F0DF89918111FCCFDF91CF91A9
24 :1001700008955C9890E4899F600D1124862F805823
25 :1001800070DF86EE90E00197F1F700000895229A63
26 :10019000559A089582E184BD84E085BD85E487BD0C
27 :1001A000089575DF60E080E0E4DF81E191E0D3DF76
28 :1001B00060E081E0DEDF82E291E0CDDF0895E7DFED
29 :1001C000E9DF0895FCDFEDDF35E423E383E090E031
30 :1001D00037BD2A984FEF5FEF6CE2415050406040CE
31 :1001E000E1F700C0000027BD2A9A4FEF5FEF6CE2F5
32 :1001F000415050406040E1F700C00000019741F7D6
33 :1002000017BC8FEF9FE34CEE815090404040E1F7E8
34 :0A0210000C00000DBCF894FFCF20
35 :10021A002020202020202020202020202020D4
36 :10022A00002D204578706572696D656E74203261A3
37 :10023A002D0020506F6C697A65692020202020CB
38 :04024A002020000070
39 :00000001FF

```

The bytes in the first line are:

Byte0	Byte1	Byte2	Byte3	Byte4	Byte5	...
\$10_{16}\$	\$00_{16}\$	\$00_{16}\$	\$00_{16}\$	\$19_{16}\$	\$\rm C0_{16}\$	...

But what is the decimal value of these numerals?

This transfer is also possible in the process shown in 2.1.1:

$Z_{10} = \sum_{i=-n}^m z_i \cdot B^i$, where z_i is the digit on position i .

As an example the Byte4 shall be written in decimal:

$Z_{10}(\text{Byte4}) = Z_{10}(19_{16}) = \sum_{i=0}^1 z_i \cdot 16^i = 1 \cdot 16^1 + 9 \cdot 16^0 = 16 + 9 = 25_{10}$

Octal

Another - much less common - base is 8. This number system is called **octal**. In this case, only $0,1,2,3,4,5,6,7$ can be used for digits.

Similar to the number systems before, the following formula applies for the transfer in decimal: $Z_8(X) = \sum_{i=-n}^m z_i \cdot 8^i$

2.1.4 From decimal to other bases

Fast Approach: Try and Subtract

Up to this point chapter, we only have converted other B -based numerals to decimal.

Now we want to try to convert a decimal into hexadecimal, for example, the numeral 315_{10} . One way is to first convert to binary by “try and subtract” and then convert the nibbles into hexadecimal:

For this example, we first try to subtract the highest power of 2, which not results in a negative number.

$$\begin{aligned} 315_{10} &- 256_{10} = 59_{10} \quad \quad 2^8 \quad \quad 59_{10} \\ &- 32_{10} = 27_{10} \quad \quad 2^5 \quad \quad 27_{10} \\ &- 16_{10} = 11_{10} \quad \quad 2^4 \quad \quad 11_{10} \\ &- 8_{10} = 3_{10} \quad \quad 2^3 \quad \quad 3_{10} \\ &- 2_{10} = 1_{10} \quad \quad 2^1 \quad \quad 1_{10} \\ &- 1_{10} = 0_{10} \quad \quad 2^0 \end{aligned}$$

The number 315_{10} is similar to $2^8 + 2^5 + 2^4 + 2^3 + 2^1 + 2^0$. This is equal to 100111011_2 or $0001\;0011\;1011_2$. For the hexadecimal value, these nibbles have to be converted, which leads to $13B_{16}$. This is often the fastest way but is based on the constraint, that one remembers the power of 2.

Works every time: Repeated Division / Multiplication

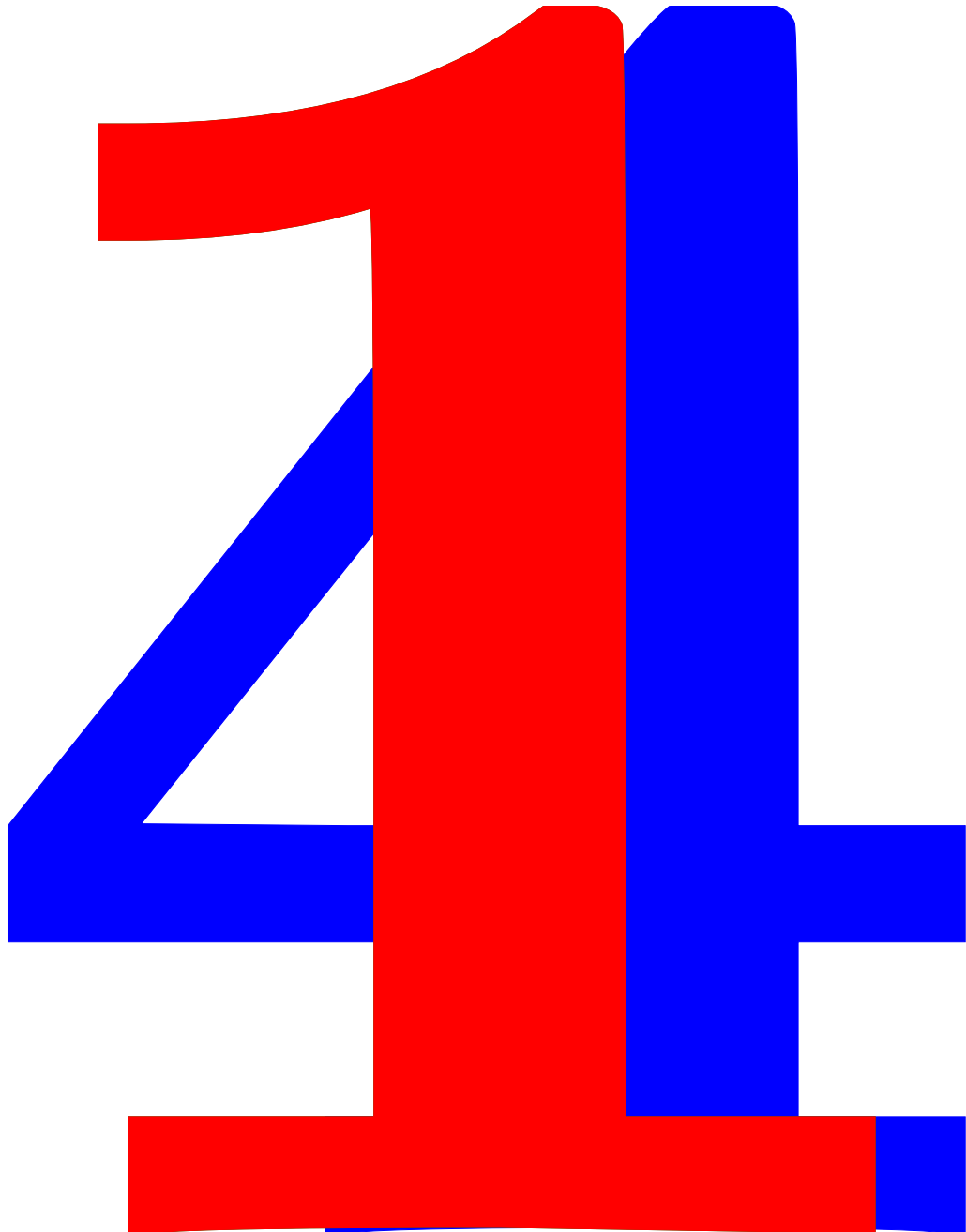
But how can we convert a decimal numeral like 452.12_{10} e.g. to hexadecimal?

The easiest way is to at first separate the value in **integer decimal z** and **decimal places f** :

$452.12 \rightarrow z = 452, f = 0.12$

1. handling the integer decimal

The integer decimal z can be converted by the abovementioned method. A different way is via repeatedly dividing by the base (here $B=16$) and using the remainder:



415C A15C A1 5C A15B A1

For each of the shown steps, the integer of the division of the step before is used (\$28\$, \$1\$).
The last steps (and any following) result in a remainder of \$0\$.
For hexadecimal numerals, we have to focus on the remainder from below to top and convert
the value into hexadecimal digits

each position in decimal	$\text{\color{red}{1}}$	$\text{\color{green}{12}}$	$\text{\color{blue}{4}}$
each position in hexadecimal	$\text{\color{red}{1}}$	$\text{\color{green}{C}}$	$\text{\color{blue}{4}}$

The result is $\text{\color{red}{1}}\text{\color{green}{C}}\text{\color{blue}{4}}_{16}$

2. handling the decimal places

The decimal places \$0.12\$ can be converted by repeatedly multiplying by the base (here \$B=16\$) and using the integer:



For each of the shown steps, the decimal places of the multiplication of the step before are used (\$\$.92 \rightarrow .72 \rightarrow .52 \rightarrow .32 \rightarrow .12 \rightarrow .92\$\$\$). The decimal place in the second last line is equal to the first one. Therefore also any further places will also be equal. That leads to repeating decimals. For hexadecimal numerals, we have to focus on the integer from top to below and convert the value into hexadecimal digits.

each position in decimal	$\textcolor{blue}{1}$	$\textcolor{green}{14}$	$\textcolor{brown}{11}$	$\textcolor{red}{8}$	$\textcolor{grey}{5}$	$\textcolor{blue}{1}$	$\textcolor{green}{14}$	...
each position in hexadecimal	$\textcolor{blue}{1}$	$\textcolor{green}{E}$	$\textcolor{brown}{B}$	$\textcolor{red}{8}$	$\textcolor{grey}{5}$	$\textcolor{blue}{1}$	$\textcolor{green}{E}$	

The result is $\textcolor{blue}{0}.\overline{\textcolor{blue}{1} \textcolor{green}{E} \textcolor{brown}{B} \textcolor{red}{8} \textcolor{grey}{5}}$

Results:

1. A decimal numeral has to be separated into integer decimal and decimal places.
2. By dividing/multiplying with the base the integer decimal / decimal places can be converted to another base. This works for all other bases like \$2\$ or \$8\$.
3. The results have to be converted (at least for the base \$B > 10\$, like hexadecimal).
4. A small decimal place can lead to longer (or even infinite) numerals in another base.

Especially the last result has a major impact on calculations on microcontrollers and computers: The internal logic is only based on binary, which also shows this problem. However, the internal memory for a numeral is limited.

Even when stored in 32bit - it is not possible to exactly convert the \$0.12_{10}\$ to binary. In the following table, the \$n\$-bit equivalent of \$0.12_{10}\$ in the binary and hexadecimal system is shown. Additionally, this value is also re-converted to a decimal numeral.

number of bits \$n\$	number system	numeral
\$8\$	binary	\$0.0001\; 1111_2\$
	hex	$\textcolor{brown}{0}.\textcolor{red}{1}\textcolor{blue}{F}_{16}$
	equiv. dec	0.12109375_{10}
\$16\$	binary	\$0.0001\; 1110\; 1011\; 1000_2\$
	hex	$\textcolor{brown}{0}.\textcolor{red}{1}\textcolor{blue}{E}\textcolor{green}{B}_{16}$
	equiv. dec	0.11999511718_{10}
\$24\$	binary	\$0.0001\; 1110\; 1011\; 1000\; 0101\; 0010_2\$
	hex	$\textcolor{brown}{0}.\textcolor{red}{1}\textcolor{blue}{E}\textcolor{green}{B}\textcolor{red}{5}\textcolor{blue}{1}_{16}$
	equiv. dec	$0.11999994516..._{10}$

number of bits n	number system	numeral
32	binary	$0.0001; 1110; 1011; 1000; 0101; 0001; 1110; 1100_2$
	hex	$0.1EB851EC_{16}$
	equiv. dec	$0.12000000011..._{10}$

This might seem like a little issue. But there are a lot of areas, where exact decimal places are mandatory, like banking, or simulations.

This situation even arises, when using [floating point representation](#) - the principal problem is also nicely explained in [this clip](#).

2.1.5 Binary Coded Decimals

The first approach to this was the development of **Binary Coded Decimals**.

The encoding algorithm from a decimal numeral into BCD is simple: *“don't use the division/multiplication method mentioned before - just take the same hexadecimal digit on each decimal position like the decimal one”*.

This means the decimal numeral 391.21_{10} is encoded to 391.21_{BCD} . Inside the processor, each digit is handled as a hexadecimal number: $3;9;1;.;2;1_{\text{BCD}}$ equals $0011;1001;0001;.;0010;0001_2$.

The main disadvantage is the ineffective storage management and more complex algorithms for addition, subtraction, and so on.

2.1.6 Marking the base of a numeral

Up to here, the marking was done by the subscript. Postfixes cannot be used in the software development environment, and sometimes are also not used in datasheets. Alternative ways for marking are the following:

base	subscripted (mathematically)	prefixed (in code)	postfixed	further marking
2 (binary)	$0010\ 1010_2$	$0b00101010$	$00101010b$	$\%00101010\$, 00101010b$
10 (decimal)	1027_{10}	1027	$1027D$	-
8 (octal)	1027_8	01027	$1027O$	-
8 (hexadecimal)	$27D_{16}$	$0x27D$	$27H$	$\$27, 27h$

Be aware, that in the code 01027 is not equal to 1027 !

When to use which base?

When programming code for an embedded system, the system will always see 0's and 1's after compiling. So, the microprocessor does not have to be considered.

In some cases, on the other hand, the binary or hexadecimal numeral is much more convenient to read:

Here is an example, where the binary numeral shows a smiley (e.g. for writing this on the screen), but the hexadecimal (or decimal) value does not give a clue what the output will be:

```
... int 0000000000000000 = { ...
```

Text is not SVG - cannot display

2.2 basic arithmetic operations in binary and hexadecimal

In this subchapter, we will have a look at the way how the arithmetic operations have to be executed manually in other bases. For math, it does not matter in which number system one calculates: a calculation like $2_{10} + 5_{10} = 7_{10}$, will be in binary $0010_2 + 0101_2 = 0111_2$. The values behind the numerals are still the same, they are only encoded differently.

The execution is similar to the well-known experience of calculating in the decimal system.

Important for all of the operations: Never forget the **carry**!

Just a small reminder in decimal:

$$\begin{array}{c} \text{\color{white}\scriptsize +192 \hspace{-0.7em} || \hspace{-0.7em} +378 \hspace{-0.7em} || \hspace{-0.7em} +672 \hspace{-0.7em} || } \\ \text{\color{red}\scriptsize 1\hspace{-0.7em} \overset{\hspace{-0.7em} \color{red}\scriptsize 2}\hspace{-0.7em} }\hspace{-0.7em} \color{red}\scriptsize 1\hspace{-0.7em} \overset{\hspace{-0.7em} \color{red}\scriptsize 1}\hspace{-0.7em} \color{red}\scriptsize 4\hspace{-0.7em} \overset{\hspace{-0.7em} \hspace{-0.7em} 2}\hspace{-0.7em} } \end{array}$$

The red numbers are the carry-over of the calculation to the right. If a calculation exceeds the limits of the base, a new digit is added and the carry is taken into account in it. A similar carry will happen in the next subchapters.

2.2.1 Addition

In Binary

The following examples shall show the concept of binary:

$$\begin{array}{|l|l|l|} \hline \text{a) } \overline{\overline{0_2}} \& \text{b) } \overline{\overline{0_2}} \& \text{c) } \overline{\overline{1_2}} \& \text{d) } \overline{\overline{1_2}} \\ \hline \end{array}$$

The four simplest examples show the carry only for $d \equiv 1 \pmod{2}$. Now we can also connect the number system with the logic gates! The addends are called A and B (e.g. $A=1_2$, $B=1_2$), the sum is the numeral CS (e.g. $C=1, S=0 \rightarrow CS=10_2$).

When analyzing the calculations above, we need one logic gate for \$\$\$, which only results in \$1\$, when either \$A=1\$ or \$B=1\$: this is the XOR gate.

For the carry \$C\$ we only get \$1\$, when both of the inputs are one: Here an AND gate has to be used.

The [figure 4](#) shows the resulting logic. This is also called a half-adder.

Fig. 4: Simulation of a half adder

The next step shall be a bit more complicated - or better: some bit more complicated. We want to do the calculation $A + B = 0011_2 + 0111_2$.

```
\begin{align*} \color{white}{+}\,00\boldsymbol{1_2} \, + 01\boldsymbol{1_2} \\ \hline \color{white}{+}\overset{\color{red}{1}}{1}\overset{\color{red}{1}}{0}\overset{\color{red}{1}}{0}\overset{\color{red}{1}}{0} \end{align*}
```

The rightmost, first step is easy since we already had this in the examples before. The next step (marked in bold) is differing: not only do we have to consider the digits from A and B , but also the carry from the calculation before. The carry from before has to be added for all the next steps similarly.

For the calculation by hand, we did four individual additions from the right to the left. The processor has to do the same. But first, we have to expand the half-adder. For a complete addition step, we need a logic with the inputs A , and B and carry from the step before C_{in} . The output will be still S and C .

Fig. 5: Simulation of a full adder

In [figure 5](#) the full adder is shown. It is an alternative representation from what we have done for calculation by hand:

1. first add one bit from A to one bit from B (half adder). The result is in S' and C' .
2. Then, take the result S' and add it with C_{in} (second half adder). The result is in S and C'' . C is already the wanted output.
3. For the carry output we need to consider both carries C' and C'' . When two or all of the inputs A , B , and C_{in} are high, then the carry has to be set to high. This can be implemented with $C' \text{ OR } C''$.

This full adder can now be stacked together to add multiple bits $A_0, A_1, \dots$ to $B_0, B_1; \dots$.

Fig. 5: Simulation of a 4-bit adder

In Hexadecimal

The calculation in hexadecimal is conceptually the same. Some examples, which we will discuss below:

```
\begin{align*} \begin{array}{lll} \text{a)} \, \color{white}{+}5_{16} \, + 3_{16} \\ \hline \color{white}{+}\overset{8_{16}}{8} \, \& \, \text{b)} \, \color{white}{+}7_{16} \, + 3_{16} \\ \hline \color{white}{+}\overset{\text{rm } A_{16}}{A} \, \& \, \text{c)} \, \color{white}{+}1_{16} \, + \text{rm } D_{16} \\ \hline \color{white}{+}\overset{\text{rm } E_{16}}{E} \, \& \, \text{d)} \, \color{white}{+}\text{rm } E_{16} \, + \text{rm } A_{16} \\ \hline \color{white}{+}\overset{\color{red}{1}}{1}\overset{8_{16}}{8} \end{array} \end{align*}
```

a) This one is simple: looks like a decimal formula..

b) Here, the summands look like decimal numerals, but the result $7_{10} + 3_{10} = 10_{10}$ is still within the range of the base. The correct symbol would be $10_{10} = \text{rm } A_{16}$

c) Now, the summands are a “bit more hexadecimally”. The easiest way is: “convert the single digit from hex to decimal, do the operation, and re-convert to hex”. For the given example: $1_{10} + 13_{10} = 14_{10} = \text{D}_{16}$

d) Also for this calculation the described way is beneficial: $\text{E}_{16} + \text{A}_{16} = 14_{10} + 10_{10} = 24_{10}$. The result is larger than the base, and therefore the value has to be separated in more digits: $24_{10} = 16_{10} + 8_{10} = 10_{16} + 8_{16} = 18_{10}$

For a hexadecimal value with more digits, the carry of the calculation before has to be added - otherwise, every step remains the same.

2.2.2 Subtraction

In Binary

The following examples shall show the concept of binary:

```
\begin{align*} \begin{array}{lll} \text{a) } 0_2 - 0_2 \\ \text{b) } 1_2 - 1_2 \\ \text{c) } 1_2 - 0_2 \\ \text{d) } 0_2 - 1_2 \end{array} \\ \begin{array}{lll} \overline{0_2} \\ \overline{0_2} \\ \overline{1_2} \\ \overline{0_2} \end{array} \\ \begin{array}{lll} \overline{0_2} \\ \overline{0_2} \\ \overline{1_2} \\ \overline{1_2} \end{array} \end{align*}
```

The calculation for d) shows the carry, which here has to borrow a bit from the next upper digit.

This is similar to the calculation: $42_{10} - 23_{10} = 19_{10}$

With more digits the calculation in the binary system will look like the following:

```
\begin{align*} 1010_2 - 011_2 \\ \overline{1010_2} \\ \overline{011_2} \end{align*}
```

In this example the bold column will be explained shortly:

```
\begin{align*} 10_2 - 1_2 \\ \overline{10_2} \\ \overline{1_2} \end{align*}
```

The calculation has to be executed as follows: $10_2 - (1_2 + 1_2) = 8_2$. Additionally, another carry has to be taken from the next digit.

In Hexadecimal

The calculation in hexadecimal is conceptually again the same. Some examples, which we will discuss below:

```
\begin{align*} \begin{array}{lll} \text{a) } 15_{16} - 13_{16} \\ \text{b) } 23_{16} - \text{B}_{16} \\ \text{c) } 3F_{16} - 1A_{16} \\ \text{d) } 38_{16} + \end{array} \\ \begin{array}{lll} \overline{15_{16}} \\ \overline{23_{16}} \\ \overline{3F_{16}} \\ \overline{38_{16}} \end{array} \end{align*}
```

$$\overline{1A_{16} + B_{16}}$$

a) This one is (again) simple: looks (again) like a decimal formula...

b) Here, both hexadecimal numerals look like decimal numerals, but the result $3_{10} - 6_{10} = 7_{10} + C_{10}$ lead to an underflow, where we have to take a carry of $C = -10_{10}$ for a decimal calculation. In hexadecimal, the carry differs. A better way to solve such a subtraction in hexadecimal is to add the carry before: $3_{16} - 6_{16} + \textcolor{red}{10}_{16}$. Then the calculation can be converted to decimal: $3_{10} - 6_{10} + \textcolor{red}{16}_{10} = 19_{10} - 6_{10} = 13_{10}$. This result has to be converted back to hexadecimal: $\text{rm } 13_{10} = D_{16}$. For the next column, the carry has to be considered.

c) For this formula the idea from b) helps, too: for each digit-by-digit subtraction, we will have a look, at whether a transformation to decimal is beneficial. For the rightmost it is: $\text{F}_{16} - A_{16} = 15_{10} - 10_{10} = 5_{10}$. The rightmost result is now: $5_{10} = 5_{16}$

d) We can do the same thing here, for $8_{16}-A_{16}$: consider the carry, convert to decimal, calculate, and re-convert to hexadecimal. $8_{10}-10_{10}+\textcolor{red}{16_{10}}=24_{10}-10_{10}=14_{10}$. The result is E_{16} .

Note!

The subtraction within the processor is done a bit differently, with the **two's complement**. In short: the range of numerals will be divided in such a way that negative numbers can also be represented. By this, the subtraction can be transformed into an addition. It is roughly like transforming $10_{10} - 4_{10}$ into $10_{10} + (-4)_{10}$.

2.2.3 Multiplication and Division

The multiplication table for binary is simple:

$$\begin{aligned} 0_2 \cdot 0_2 &= 0_2 \quad \parallel \quad 0_2 \cdot 1_2 = 0_2 \quad \parallel \quad 1_2 \cdot 0_2 = 0_2 \quad \parallel \quad 1_2 \cdot 1_2 = 1_2 \\ \end{aligned}$$

Multiplication of longer binary numerals has to be executed similarly to decimal numerals:

[illegible]

The division is also relatively simple:

$$\begin{array}{l} 0 \ 2 : 1 \ 2 = 0 \ 2 \\ 1 \ 2 : 1 \ 2 = 1 \ 2 \end{array}$$

With longer numerals it looks like the following:

Fig. 7: binary division $10101_2 : 11_2 = 111_2$

related Links

- [Conversion tool from decimal to hexadecimal](#), this tool shows also the steps and can be used vice versa
- A short video on the “[Everything Formula](#)” which prints all possible output and is directly related to binary encoding.

Exercises

Exercise 2.3.1. power tables

fill out the tables and remember the numerals!

Power...	Power...	Decimal numeral			
2^{13}	—				
2^{12}	16^3	4	0	9	6
2^{11}	—				
2^{10}	—				
2^9	—				
2^8	16^2				
2^7	—				
2^6	—				
2^5	—				
2^4	16^1				
2^3	—				
2^2	—				
2^1	—				
2^0	16^0				1

Fig. 8: Important numerals 1

Fig. 9: Important numerals 2

Power...	Power...	Decimal numeral							
2^{-1}	—	0.	5						
2^{-2}	—	0.							
2^{-3}	—	0.							
2^{-4}	16^{-1}	0.							
2^{-5}	—	0.							
2^{-6}	—	0.							
2^{-7}	—	0.	0	0	7	8	1	2	5

Text is not SVG - cannot display

Binary				Decimal	Hex	BCD
0	0	0	0			
1	1	1	1			

Fig. 10: Important numerals 3

Exercise 2.3.2. Conversion to Binary

Convert the following hexadecimal numerals to the binary system.
Tip: Convert nibble by nibble to binary.

- 1. 4321H
- 2. 0xCAFE
- 3. 4be79
- 4. 398B4C
- 5. 0x3a4f.4ecd
- 6. 1.3DAF

Exercise 2.3.3. Conversion to Hexadecimal and Decimal

Convert the following numerals from dual to hexadecimal and decimal numerals.

Tip: Use the numerals from Exercise 2.3.1.

1. 0b10110110
2. $\text{\%10100101}$
3. $\text{1111}, \text{0111}, \text{0001}, \text{0011B}$
4. 0110111.0101b
5. 0.10101b
6. $\text{0.1001}, \text{1001}, \dots \text{b}$

Exercise 2.3.4. Conversion to Hexadecimal and Binary

Convert the following numerals from decimal to hexadecimal and dual numerals.

Tip: Use the numerals from Exercise 2.3.1.

1. 123_{10}
2. 637_{10}
3. 1777_{10}
4. 8999_{10}
5. 41.4_{10}
6. 0.9_{10}
7. 11.3_{10}

Exercise 2.3.5. Addition and Subtraction

Execute the following operations manually with the given dual numerals.

1. $\text{01110111B} + \text{11010101B}$
2. $\text{01011001B} + \text{10011111B}$
3. $\text{01100111B} - \text{01001100B}$
4. $\text{10101110B} - \text{10000111B}$
5. $\text{0011.111B} - \text{0011.100B}$
6. $\text{001111.011B} + \text{1.010011B}$
7. $\text{01101.100B} - \text{01000.111B}$
8. $\text{100101.11B} - \text{110001.10B}$

Exercise 2.3.6. Multiplication and Division

Execute the following operations manually with the given dual numerals.

1. $\text{1011B} \cdot \text{0101B}$
2. $\text{01100101B} \cdot \text{110B}$
3. $\text{10101111B} : \text{101B}$
4. $\text{11110000B} : \text{1000B}$

Exercise 2.3.7. Operation in Hex

Execute the following operations manually with the given dual numerals.

1. $\text{\rm 1334H} + \text{\rm 07ABH}$
2. $\text{\rm 0DC43H} - \text{\rm 0BD19H}$
3. $\text{\rm 1F23H} + \text{\rm 90E8H}$
4. $\text{\rm 98C5H} - \text{\rm 84CAH}$
5. $\text{\rm 234AH} + \text{\rm 7EE6H}$
6. $\text{\rm 0F10CH} - \text{\rm 0ED43H}$

Exercise 2.3.8. Further Questions

1. Write down the decimal numeral for the first 4 hexadecimal places
2. Which possible value range can be covered with an 8-bit variable?

Exercise 2.3.9. One's Complement

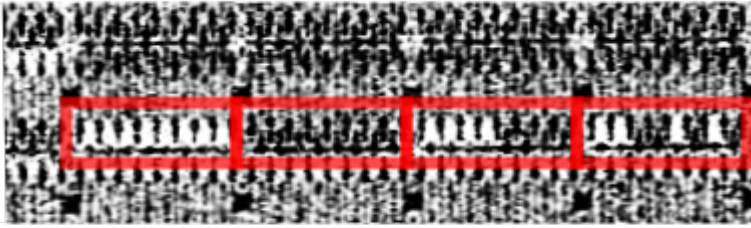
In the simulation in [figure 11](#) the one's complement of a 4bit value is shown. On initialization the value is the following:

- The value for A is $1010_2 = 10_{10}$. The decimal value 10 is also shown in the first display from the left in the circuit.
- The value for B is $1110_2 = 14_{10}$. The decimal value 14 is also shown in the second display from the left in the circuit.
- The addition $A+B$ leads to $S = 1000_2 = 8_{10}$. The decimal value 10 is also shown in the rightmost display in the circuit.
- There is another value in the second display from the right. This is called B' and it is equal to $0010_2 = 2_{10}$.

Fig. 11: One's Complement

1. Why is $10 + 14 = 8$? What happens to other values?
The input values A and B can be changed by clicking on the bit values.
2. Try to analyze how B' (shown in the brackets) is derived from B .
3. Have a look at the wiki page of the [Ones' complement](#) to understand how negative numbers are represented in a microcontroller.
4. Imagine that you have to rescue data from an old storage device. The interesting bits are given by the boxed area in [figure 12](#). The bytes are LSB 0 oriented.
 1. What are these values in an unsigned integer?
 2. What are these values in signed integers based on the ones' complement?

Fig. 12: example for stored data in an EEPROM



From: Courbon F., Skorobogatov S., Woods C. (2017) Reverse Engineering Flash EEPROM Memories Using Scanning Electron Microscopy. In: Lemke-Rust K., Tunstall M. (eds) Smart Card Research and Advanced Applications. CARDIS 2016. Lecture Notes in Computer Science, vol 10146. Springer, Cham. https://doi.org/10.1007/978-3-319-54669-8_4

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